

CONTROLLED RANDOM WALK WITH A TARGET SITE

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ABSTRACT. We consider a simple random walk $\{W_i\}$ on \mathbb{Z}^d , $d = 1, 2$, in which the walker may choose to stand still for a limited time. The time horizon is n , the maximum consecutive time steps which can be spent standing still is m_n and the goal is to maximize $P(W_n = 0)$. We show that for $d = 1$, if $m_n \gg (\log n)^{2+\gamma}$ for some $\gamma > 0$, there is a strategy for each n yielding $P(W_n = 0) \rightarrow 1$. For $d = 2$, if $m_n \gg n^\epsilon$ for some $\epsilon > 0$ then there are strategies yielding $\liminf_n P(W_n = 0) > 0$.

1. INTRODUCTION.

We consider a process $\{W_i, 0 \leq i \leq n\}$ on \mathbb{Z}^d ($d = 1$ or 2) in which $W_0 = 0$ and each step $W_i - W_{i-1}$ either is 0 (i.e. standing still) or is a step $\pm e_i$ of symmetric simple random walk (SSRW), with e_i being the i th unit coordinate vector. The choice between standing still and SSRW step is determined by a strategy. Formally, a *strategy* (or *n-strategy*) is a mapping $\delta_n : \mathcal{T}_n \rightarrow \{0, 1\}$ defined on the space

$$\mathcal{T}_n = \{(i, w_i), 0 \leq i \leq j\} : 0 \leq j \leq n-1, |w_i - w_{i-1}| \leq 1 \text{ for all } i\}$$

of all space-time trajectories of length less than n ; here the values 0 and 1 for δ_n correspond to standing still and taking a SSRW step, respectively. Thus the choice of whether to take a step at a time i depends on the trajectory up to time $i-1$. Letting $\{Y_i, i \geq 1\}$ be SSRW, we then construct the controlled random walk process iteratively from the strategy by

$$W_0 = 0, \quad W_i - W_{i-1} = \Delta_i(Y_i - Y_{i-1}) \quad i \leq n,$$

where

$$\Delta_i = \delta_n \left(\{(k, W_k) : 0 \leq k \leq i-1\} \right).$$

We define the time since the last SSRW step to be

$$J_i = \min\{j \in [0, i-1] : W_{i-j} \neq W_{i-j-1}\},$$

with $J_i = i$ if the set in this definition is empty. We designate a maximum number $m_n - 1$ of consecutive steps standing still and say that a strategy is *admissible* if $J_i \leq m_n - 1$ a.s. for all $i \leq n$. A strategy is *Markov* if the value of $\delta_n(\{(k, W_k) : 0 \leq k \leq i-1\})$ depends only on the pair (W_i, J_i) . It is easily seen that if the strategy is Markov, then $\{(W_i, J_i) : i \geq 0\}$ is a Markov process. We write \mathcal{A}_n for the set of all admissible strategies and \mathcal{A}_n^* for the set of all admissible Markov strategies. When the dependence on the

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strategy needs to be made clear, we write $P(\cdot; \delta_n)$ for probability when strategy δ_n is used, but generally we suppress the δ_n in the notation.

We are interested in steering the process toward a target site, and specifically in the behavior of

$$(1.1) \quad \hat{P}(W_n = 0) = \sup_{\delta_n \in \mathcal{A}_n} P(W_n = 0; \delta_n)$$

as $n \rightarrow \infty$. The word “steering” is a bit misleading here, as the process never has a drift. For SSRW, of course $P(W_n = 0) \sim Cn^{-d/2}$. (Here and throughout the paper, C and C_1, C_2, \dots are generic constants, and $a_n \sim b_n$ means the ratio converges to 1.) A simple strategy to increase this is to minimize steps, i.e. always stand still as long as allowed, taking only n/m_n steps, yielding $P(W_n = 0) \asymp (m_n/n)^{d/2}$, where we use $a_n \asymp b_n$ to mean the ratio is bounded away from 0 and ∞ . A slightly more sophisticated strategy is to minimize steps until time $n - m_n$ then maximize steps (i.e. never stand still) until $\{W_i\}$ hits 0 (if it does), and then stand still until time n . For $d = 1$, conditionally on $\{|W_{n-m_n}| \leq m_n^{1/2}\}$ this strategy has a success probability of order 1, so the unconditional success probability is of order $\min(m_n/n^{1/2}, 1)$; in particular it is bounded away from 0 if $m_n \geq Cn^{1/2}$.

This brings up two questions. For which $\{m_n\}$ is $\hat{P}(W_n = 0)$ bounded away from 0? And are there nontrivial $\{m_n\}$ for which $\hat{P}(W_n = 0) \rightarrow 1$? Our two main theorems give some answers.

Theorem 1.1. *Suppose $d = 1$ and $m_n \gg (\log n)^{2+\gamma}$ for some $\gamma > 0$. Then $\hat{P}(W_n = 0) \rightarrow 1$ as $n \rightarrow \infty$.*

In one dimension, the “more sophisticated” strategy described above relies on the fact that a SSRW started at distance k from 0 has a probability of order one to hit 0 by time k^2 . In two dimensions, this probability is only of order $(\log k)^{-1}$ so it is more difficult to construct a strategy based on waiting for the RW to return to 0 after it has wandered away. Nonetheless we have the following.

Theorem 1.2. *Suppose $d = 2$ and $m_n \gg n^\epsilon$, for some $\epsilon > 0$. Then*

$$\liminf_n \hat{P}(W_n = 0) > 0.$$

In (1.1) one could consider only Markov strategies, i.e. take the sup over \mathcal{A}_n^* . Since the underlying SSRW is Markov, one cannot actually do better with a non-Markov strategy, so the two sups are the same. Allowing non-Markov strategies simply lets us use ones that can be concisely described and analyzed.

Remark 1.3. One could consider alternate ways of slowing down the controlled RW, in place of standing still. For example, one could allow a choice between a standard SSRW step and a delayed SSRW step, the latter meaning we stand still with probability $1 - \frac{1}{m_n}$ and take a standard SSRW step of $\pm e_i$ with probability $1/m_n$, with no limit on how many consecutive times we choose the delayed SSRW step. A brief examination of the proofs shows that both theorem statements above remain valid in this case.

Remark 1.4. The continuous analog of the problem we consider is a diffusion ξ_t in which the drift is always 0 and one can control the diffusivity $\sigma(x, t)$, but constrained to an

interval $[\sigma_1, \sigma_2]$, on a time interval $[0, T]$, where $\sigma_1 > 0$. We take $[\sigma_1, \sigma_2] = [\epsilon_T, 1]$ and ask, how slowly can we have $\epsilon_T \rightarrow 0$ as $T \rightarrow \infty$ and still have $\hat{P}(|\xi_T| < 1) \rightarrow 1$ or $\liminf_T P(|\xi_T| < 1) > 0$? Note that ϵ_T is the analog of $1/m_n$. McNamara [3] considered closely related questions for a one-dimensional diffusion; see also [2] for another variant. He proved that there are constants $K = K(\sigma_1, \sigma_2)$ and $\beta = \beta(\sigma_1/\sigma_2)$ such that

$$(1.2) \quad \hat{P}^x \left(|\xi_T| \leq \frac{\delta}{2} \right) \leq K \left(\frac{\delta}{\sqrt{T}} \right)^\beta \quad \text{for all } x \in \mathbb{R}, \delta > 0,$$

while in the other direction, for each $h > 0$ there exists $\Delta > 0$ such that

$$(1.3) \quad \hat{P}^0 \left(|\xi_T| \leq \frac{\delta}{2} \right) > K \left(\frac{\delta}{\sqrt{T}} \right)^{\beta+h} \quad \text{for all } \delta \in (0, \Delta).$$

Here \hat{P}^x denotes probability (maximized over the allowed controls) for a process started at x . Numerical evidence was given that β is approximately proportional to σ_1/σ_2 . If we assume this to be true and consider ϵ_T of order $1/\log T$, we get β of order $1/\log T$ as well. If we could take h also of this order, then the right side of (1.3) would be bounded away from 0 in T , as desired. The problem is that one cannot get a useful result from (1.3) if one takes h depending on T , since then δ must also depend on T . Nonetheless we may observe that $\epsilon_T \leq C/\log T$ corresponds to $m_n \geq C \log n$ in our discrete problem. Further, still assuming β proportional to σ_1/σ_2 , we see that since we want the left side of (1.2) bounded away from 0, we cannot allow $\epsilon_T \gg 1/\log T$, suggesting that we perhaps cannot do better than requiring $m_n \geq C \log n$ in Theorem 1.1.

2. PROOF OF THEOREM 1.1

Throughout the paper we will make use of various quantities which approach infinity as $n \rightarrow \infty$. For $\lambda > 0$ let $u_n(\lambda) = \max\{k : m_n^{1+k\lambda} \leq n\}$; note that $u_n(\lambda)$ is of order $\log n / \log m_n$ for all $\lambda > 0$. Also, the hypothesis that $m_n \gg (\log n)^{2+\gamma}$ for some $\gamma > 0$ is equivalent to

$$(2.1) \quad m_n \gg \left(\frac{(\log n)^2}{\log \log n} \right)^{1/(1-\eta)} \quad \text{for some } \eta \in (0, 1).$$

Fixing such an η and writing u_n for $u_n(\eta)$, we observe that (2.1) is equivalent to

$$m_n^{1-\eta} \log m_n \gg (\log n)^2,$$

hence also to

$$m_n^{1+\eta\theta(m_n^{1-\eta}/\log m_n)^{1/2}} \geq n \quad \text{for all large } n, \text{ for every } \theta > 0,$$

and therefore finally to

$$(2.2) \quad u_n = o \left(\left(\frac{m_n^{1-\eta}}{\log m_n} \right)^{1/2} \right).$$

Fix n , write m for m_n . and define

$$\epsilon_m = \frac{\log m}{m^{1-\eta}}, \quad \text{so} \quad \epsilon_m = o \left(\frac{1}{u_n^2} \right) \quad \text{as } n \rightarrow \infty.$$

We define a sequence of windows $\{t_k\} \times I_k, 1 \leq k \leq u_n$, in space-time, with size decreasing as k increases and the target $(n, 0)$ is approached; we then construct a strategy which makes the space-time trajectory of the process pass through all of these windows, with high probability. Specifically, let

$$t_k = n - m^{1+\eta(u_n-k)}, \quad 2 \leq k \leq u_n, \quad \text{and } t_0 = 0, \quad t_1 = \frac{t_2}{2}, \quad t_{u_n+1} = n,$$

$$N_k = t_k - t_{k-1}, \quad h_{mk} = (\epsilon_m N_{k+1})^{1/2}.$$

Let $I_0 = \{0\}$,

$$I_k = [-h_{mk}, h_{mk}], \quad 1 \leq k \leq u_n, \quad \text{and } I_{u_n+1} = \{0\}.$$

We want to find a strategy, and choice of ϵ_m , for which we can show

$$(2.3) \quad \max_{k \leq u_n+1} P(W_{t_k} \notin I_k \mid W_{t_{k-1}} \in I_{k-1}) = o\left(\frac{1}{u_n}\right) \quad \text{as } n \rightarrow \infty.$$

From (2.2) and (2.3) it follows that

$$(2.4) \quad P(W_n = 0) \geq \prod_{k=1}^{u_n+1} P(W_{t_k} \in I_k \mid W_{t_{k-1}} \in I_{k-1}) \geq \left(1 - o\left(\frac{1}{u_n}\right)\right)^{u_n+1} = 1 - o(1)$$

as $n \rightarrow \infty$, proving the theorem.

For (2.3), fix $1 \leq k \leq u_n + 1$ and let $\tau_0 = \tau_0^{(k)} = \min\{t > t_{k-1} : W_t = 0\}$. Fix $x \in I_{k-1}$ and observe that

$$(2.5) \quad \begin{aligned} P(W_{t_k} \notin I_k \mid W_{t_{k-1}} = x) &\leq P(\tau_0 > t_k \mid W_{t_{k-1}} = x) \\ &\quad + P(\tau_0 \leq t_k, |W_{t_k}| > h_{mk} \mid W_{t_{k-1}} = x). \end{aligned}$$

To fully specify, and then bound, these probabilities we need to designate a strategy; we do so by describing how the strategy works between t_{k-1} and t_k , for general k . For $k \leq u_n$, we begin by taking all SSRW steps (i.e. no standing still) from time t_{k-1} until time $\tau_0 \wedge t_k$. If $\tau_0 > t_k$, we deem the strategy to have failed and we continue in an arbitrary manner, say all SSRW steps. If $\tau_0 \leq t_k$, we continue from time τ_0 to t_k by always standing still for the maximum allowed period of time m during the interval $(\tau_0, t_k]$, that is, we take an SSRW step every m th time step. The last standing period is truncated if it would otherwise go beyond time t_k .

For $k = u_n + 1$, our strategy during $(t_{u_n}, t_{u_n+1}] = (n - m, n]$ is to maximize steps until the time $\tau_0 = \tau_0^{(u_n+1)}$ when the process first hits 0 (if $\tau_0 \leq n$), then stand still until time n .

We now bound the first term on the right side of (2.5). From the Reflection Principle we have (for ℓ, h_{mk} of opposite even-odd parity):

$$(2.6) \quad P(\tau_0^{(k)} - t_{k-1} > \ell \mid W_{t_{k-1}} = x) = P(-|x| < Y_\ell < |x|),$$

and hence for every $\epsilon > 0$,

$$(2.7) \quad \begin{aligned} P(\tau_0^{(k)} - t_{k-1} > yx^2 \mid W_{t_{k-1}} = x) &\sim 2 \left(\Phi\left(\frac{1}{\sqrt{y}}\right) - \Phi(0) \right) \\ &\quad \text{as } |x| \rightarrow \infty, \text{ uniformly over } y \in [\epsilon, \infty). \end{aligned}$$

The left side of (2.6) is a nondecreasing function of $|x|$, so for all $x \in I_k$, by (2.7)

$$(2.8) \quad P(\tau_0 > t_k \mid W_{t_{k-1}} = x) \leq P(\tau_0 - t_{k-1} > N_k \mid W_{t_{k-1}} = h_{mk}) \leq C\sqrt{\epsilon_m}.$$

Turning to the second term on the right side of (2.5), it is 0 for $k = u_n + 1$ so we consider $k \leq u_n$. We can condition also on τ_0 as follows: for $t \in (t_{k-1}, t_k]$, using Hoeffding's Inequality [1],

$$(2.9) \quad \begin{aligned} P(|W_{t_k}| > h_{mk} \mid \tau_0 = t, W_{t_{k-1}} = x) &= P(|W_{t_k}| > h_{mk} \mid W_t = 0) \\ &= P(|Y_{(t_k-t)/m}| > h_{mk}) \\ &\leq e^{-h_{mk}^2 m / 2N_k} \\ &= e^{-m\epsilon_m N_{k+1} / 2N_k}. \end{aligned}$$

For $k \geq 1$ we have $N_{k+1}/N_k \geq m^{-\eta}$ (with equality for $k \geq 3$), so (2.9) says that

$$(2.10) \quad P(|W_{t_k}| > h_{mk} \mid \tau_0 = t, W_{t_{k-1}} = x) \leq m^{-1/2} \leq \sqrt{\epsilon_m}.$$

Since $x \in I_{k-1}$ is arbitrary, combining (2.5), (2.8) and (2.10) yields that for all $1 \leq k \leq u_n + 1$,

$$(2.11) \quad P(W_{t_k} \notin I_k \mid W_{t_{k-1}} \in I_{k-1}) \leq C\sqrt{\epsilon_m} = o\left(\frac{1}{u_n}\right) \quad \text{as } n \rightarrow \infty,$$

proving (2.3) and thereby proving Theorem 1.1.

3. PROOF OF THEOREM 1.2

We keep the same definition of $u_n(\lambda)$ and note that now our hypothesis on m_n is equivalent to the statement that $\{u_n\}$ is bounded, say $u_n \leq u < \infty$ for all n . We keep the same formula for t_k but with η replaced by κ , determined as follows. Choose

$$\theta \in \left(\frac{1-\epsilon}{2}, \frac{1}{2}\right),$$

so $1 - 2\theta < \epsilon$, then choose κ small enough so

$$(3.1) \quad 0 < \frac{1 - 2\theta}{1 - 2\kappa\theta} < \epsilon.$$

We then write u_n for $u_n(\kappa)$. For our windows, in place of the interval I_k we have the square $Q_k = [-N_{k+1}^\theta, N_{k+1}^\theta]^2$. To distinguish dimensions clearly, we now write $\{Y_i^{(d)}\}$ for d -dimensional SSRW, $d = 1, 2$. We use the same strategy as in one dimension: in each interval $(t_{k-1}, t_k]$, take an SSRW step every time step until time $\tau_0^{(k)}$, then an SSRW step every m th time step from time $\tau_0^{(k)}$ to t_k . In place of (2.3), we will need that for some $C > 0$,

$$(3.2) \quad P(W_{t_k} \in Q_k \mid W_{t_{k-1}} \in Q_{k-1}) \geq C \quad \text{for all } k, \text{ for all large } n.$$

Let $L_{[a,b]}$ be the number of visits to 0 by the SSRW $\{Y_i^{(2)}\}$ during the time interval $[a, b]$. In comparison to (2.8), we have for all $x \in Q_{k-1}$:

$$(3.3) \quad \begin{aligned} P(\tau_0 \leq t_k \mid W_{t_{k-1}} = x) &= P(\tau_0 \leq t_k \mid Y_{t_{k-1}}^{(2)} = x) \\ &\geq \frac{E(L_{(t_{k-1}, t_k]} \mid Y_{t_{k-1}}^{(2)} = x)}{E(L_{(t_{k-1}, t_k]} \mid \tau_0 \leq t_k, Y_{t_{k-1}}^{(2)} = x)} \end{aligned}$$

Since $|x| \leq CN_k^\theta$ with $\theta < \frac{1}{2}$, the numerator in (3.3) is equal to

$$(3.4) \quad \sum_{i=1}^{N_k} P(Y_i = 0 \mid Y_0^{(2)} = x) \geq C \sum_{i=|x|^2}^{N_k} \frac{1}{i} \geq C_1 \log N_k.$$

Here C_1 depends on θ . The denominator in (3.3) is bounded above by

$$(3.5) \quad E(L_{(t_{k-1}, t_k]} \mid Y_0^{(2)} = 0) \leq C \sum_{i=1}^{N_k} \frac{1}{i} \leq C_2 \log N_k.$$

Therefore we have the analog of (2.8):

$$(3.6) \quad P(\tau_0 \leq t_k \mid W_{t_{k-1}} = x) \geq \frac{C_1}{C_2} > 0 \quad \text{for all } k \leq u_n, x \in Q_{k-1}.$$

In comparison to (2.9), we have for $x \in Q_k$ and $t \in (t_{k-1}, t_k]$, using again Hoeffding's Inequality [1]:

$$(3.7) \quad \begin{aligned} P(W_{t_k} \notin Q_k \mid \tau_0 = t, W_{t_{k-1}} = x) &= P(W_{t_k} \notin Q_k \mid W_t = 0) \\ &= P(Y_{(t_k-t)/m} \notin Q_k \mid Y_0^{(2)} = 0) \\ &\leq 2 \max_{j \leq N_k/m} P(|Y_j^{(1)}| > N_{k+1}^\theta) \\ &\leq 4e^{-mN_{k+1}^{2\theta}/2N_k}. \end{aligned}$$

For $k \geq 1$ we have $N_{k+1} \geq m^{-\kappa} N_k$ (with equality for $k \geq 3$), so by (3.1),

$$(3.8) \quad \frac{mN_{k+1}^{2\theta}}{N_k} \geq \frac{m^{1-2\kappa\theta}}{N_k^{1-2\theta}} \geq \frac{m^{1-2\kappa\theta}}{n^{1-2\theta}} \rightarrow \infty \text{ as } n \rightarrow \infty.$$

Combining (3.6), (3.7) and (3.8) we see that (3.2) holds. Since $\{u_n\}$ is bounded, it follows as in (2.4) that $\liminf_n P(W_n = 0) > 0$.

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